

$$\begin{aligned}
k(r)=2: & [A, B_i^r, y_i^r] \bar{y}_i^r \rightarrow [A, y_i^r] D_{\ell(r)}^r \\
& [A, B_i^r, y_i^r] [y_i^r, \$] \rightarrow [A, y_i^r] [D_{\ell(r)}^r, \$] \\
& [B_i^r, y_i^r] \bar{y}_i^r \rightarrow \bar{y}_i^r D_{\ell(r)}^r \\
& [B_i^r, y_i^r] [y_i^r, \$] \rightarrow \bar{y}_i^r [D_{\ell(r)}^r, \$] \\
k(r) > 2 & [A, B_i^r, y_i^r] \bar{y}_i^r \rightarrow [A, y_i^r] [B_{i+1}^r, y_{i+1}^r] \\
& [B_i^r, y_i^r] \bar{y}_{i+1}^r \rightarrow \bar{y}_i^r [B_{i+1}^r, y_{i+1}^r] \quad (1 \leq i < k(r)-1) \\
& [B_{k(r)-1}^r, y_{k(r)-1}^r] \bar{y}_{k(r)}^r \rightarrow \bar{y}_{k(r)-1}^r D_{\ell(r)}^r \\
& [B_{k(r)-1}^r, y_{k(r)-1}^r] [y_{k(r)}^r, \$] \rightarrow \bar{y}_{k(r)-1}^r [D_{\ell(r)}^r, \$] \\
& [D_{\ell(r)}^r, \$] \rightarrow D_{\ell(r)-1}^r [z_{\ell(r)}^r, \$] \quad (k(r) < \ell(r)) \\
& \bar{y}_{\ell(r)-1}^r [D_{\ell(r)}^r, \$] \rightarrow D_{\ell(r)-1}^r [z_{\ell(r)}^r, \$] \quad (k(r) = \ell(r)) \\
& [A, y_i^r] [D_i^r, \$] \rightarrow [\phi, z_i^r] [z_i^r, \$] \quad (k(r) = \ell(r) = 2)
\end{aligned}$$

$$\begin{aligned}
D_j^r & \rightarrow D_{j-1}^r \bar{z}_j^r \quad (k(r) < j \leq \ell(r)) \\
\bar{y}_{j-1}^r D_j^r & \rightarrow D_{j-1}^r \bar{z}_j^r \quad (2 < j \leq k(r)) \\
\bar{y}_i^r D_i^r & \rightarrow [E, z_i^r] \bar{z}_i^r \\
[A, y_i^r] D_i^r & \rightarrow [\phi, z_i^r] \bar{z}_i^r
\end{aligned}$$

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$$\begin{aligned}
[B_i^r, y_i^r] \bar{x} & \rightarrow [E, y_i^r] \bar{x} \quad (x \neq y_{i+1}^r) \\
[B_i^r, y_i^r] [x, \$] & \rightarrow \bar{y}_i^r [E, x, \$] \quad (x \neq y_{i+1}^r) \text{ ODER} \\
& (i < k(r)-1) \\
[B_i^r, x, \$] & \rightarrow [E, x, \$]
\end{aligned}$$

(4) HERSTELLEN VON  $w \in T^*$

$$\begin{aligned}
[\phi, x, \$] & \rightarrow x \quad (x \in T) \\
[\phi, x] [y, \$] & \rightarrow xy \quad (x, y \in T) \\
[\phi, x] \bar{y} & \rightarrow [A, x] [F, y] \quad " \\
[F, x] \bar{y} & \rightarrow \bar{x} [F, y] \quad " \\
[F, x] [y, \$] & \rightarrow \bar{x} [G, y] \quad " \\
\bar{x} [G, y] & \rightarrow [G, x] y \\
[A, x] [G, y] & \rightarrow xy \\
[F, x] \bar{y} & \rightarrow [H, x] \bar{y} \quad (y \notin T) \\
[F, x] [y, \$] & \rightarrow [H, x] [y, \$] \quad " \\
\bar{x} [H, y] & \rightarrow [H, x] \bar{y} \\
[A, x] [H, y] & \rightarrow [\phi, x] \bar{y}
\end{aligned}$$

DANN GILT  $L(G') = L(G)$

TH ZU JEDER MONOTONEN GRAMMATIK  $G = (V, T, P, S)$  EXISTIERT EFFEKTIV EINE MONOTONE GRAMMATIK  $G' = (V', T, P', S')$  MIT REGELN NUR DER GESTALT  $A \rightarrow BC, A \rightarrow a, AB \rightarrow CD \quad (A, B, C, D \in V-T, a \in T)$  UND EVENTUELL  $S' \rightarrow \lambda$  (DANN  $S'$  NIE RECHTS).

BEW.: NACH VORIGEM SATZ HABEN DIE REGELN DIE GESTALT  $A \rightarrow BC, A \rightarrow B, A \rightarrow a, AB \rightarrow CD, S \rightarrow \lambda$  ZU BESEITIGEN SIND:  $A \rightarrow B$