

## Wissensrepräsentation

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## Wissensrepräsentation

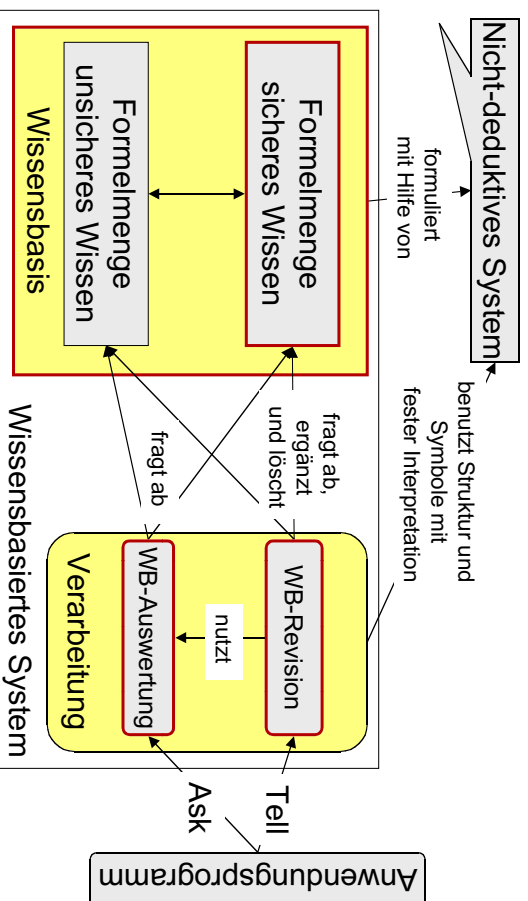
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### Sitzung 14: Vererbung / Inheritance

- Vererbung
- Pfadbasiertes Schliessen
- Vererbungsnetze

### Wissensbasiertes System mit nicht-deduktiver Wissensbasis



### Literatur

D. Gabbay, C. Hogger, and J. Robinson (eds.) (1994). *Handbook of Logic in Artificial Intelligence and Logic Programming, Volume 3: Nonmonotonic Reasoning and Uncertain Reasoning*, Oxford University Press.

#### Inbesondere

J. Horty. Some direct theories of nonmonotonic inheritance. In D. Gabbay, C. Hogger, and J. Robinson (eds.), Oxford University Press (1994), pp. 111 - 187.

D. Touretzky. (1986). The mathematics of inheritance systems. Morgan Kaufmann, Los Altos, Ca., 1986.

L. A. Stein (1992). Resolving Ambiguity in Nonmonotonic Inheritance Hierarchies, *Artificial Intelligence* 55 (2-3): 259-310.

## Vererbung (Inheritance)

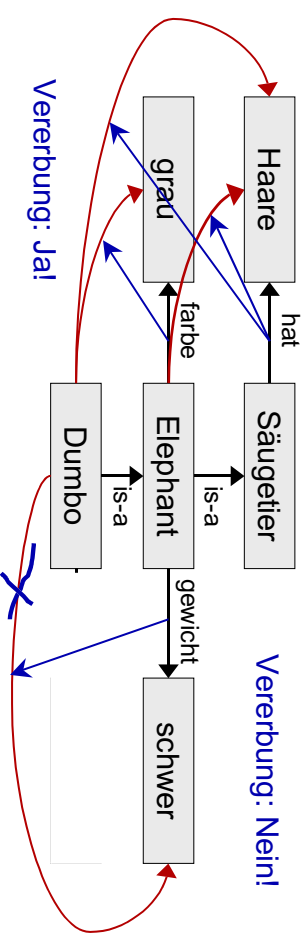
When thinking object-centered, hierarchy is a natural way to view the world

- importance of abstraction in remembering and reasoning
- groups of things share properties in the world (e.g., mammals, seafood)
- don't have to repeat representations (e.g., sufficient to say that "elephants are mammals" to know a lot about them)

z.B. Frames, Semantische Netze, Beschreibungslogik (description logics)

## Semantische Netze

- Repräsentationen für Konzept-Wissen
- Graphen mit Knoten für Konzepte und Individuen und beschriftete Pfeilen für binäre Relationen
- Definiert über Bilder, keine klare Semantik, keine spezifizierten Verarbeitungsmechanismen



Vererbung: Ja!

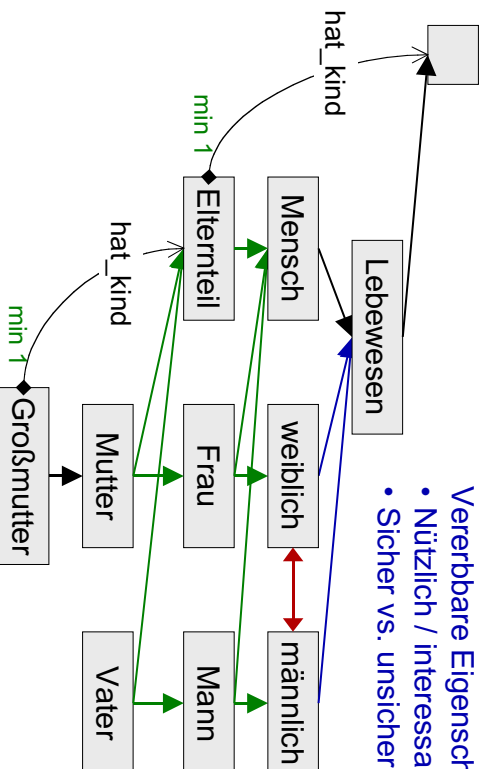
Vererbung: Nein!

## Beispiel: Konzeptsysteme

## Gruppen- diskussion

Vererbare Eigenschaften:

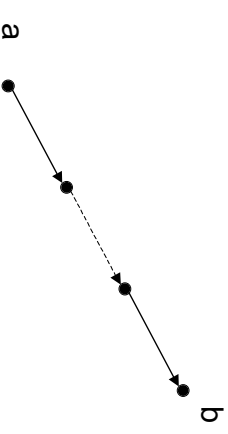
- Nützlich / interessant
- Sicher vs. unsicher



## Inheritance

“Inheritance” is the result of transitivity reasoning over paths

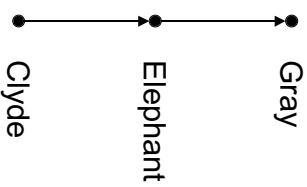
- for strict networks, *modus ponens* in graphical form
- does *a* inherit from *b*?  
≡ is *b* in the transitive closure of IS-A (or subsumption) from *a*?



## Pfad-basiertes Schliessen

Focus just on inheritance and transitivity

- many interesting considerations in looking just at where information comes from in a network representation
- abstract descriptions, and properties into *nodes* in graphs, and just look at reasoning with paths and the conclusions they lead us to



## Pfad-basiertes Schliessen (cont'd)

• edges in the network:

Clyde—Elephant, Elephant—Gray

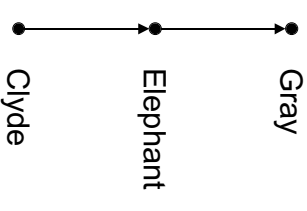
- paths included in this network: {Clyde—Elephant—Gray}, plus edges

- in general, a path is a sequence of 1 or more edges

• conclusions supported by the paths:

Clyde → Elephant; Elephant → Gray;

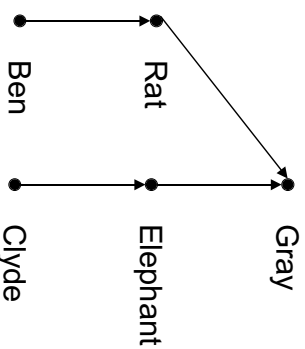
Clyde → Gray



## Vererbungsnetze (1)

**Strict inheritance in trees**

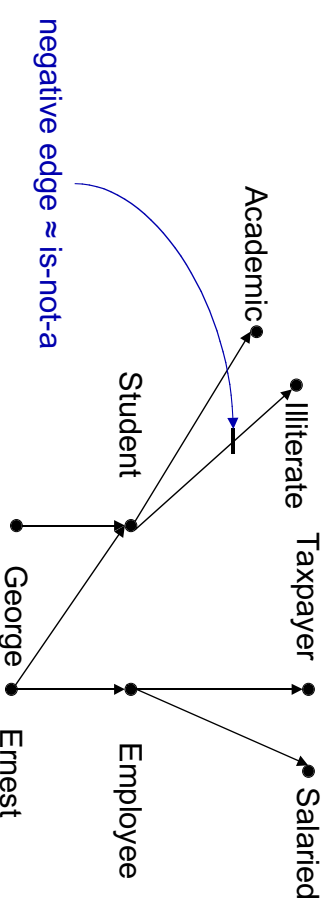
- as in description logics
- conclusions produced by complete transitive closure on all paths (any traversal procedure will do); all reachable nodes are implied



## Vererbungsnetze (2)

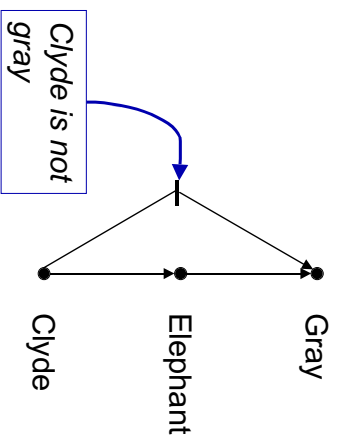
**Strict inheritance in DAGs**

- as in DL's with multiple AND parents
- same as above: all conclusions you can reach by any paths are supported



## Defeasible inheritance (1)

- inherited properties do not always hold, and can be **overridden** (defeated)
- conclusions determined by searching upward from “**focus node**” and selecting first version of property you want



- **defeasible**
- **anfechtbar**
- **annullierbar**
- „Polarität“ von Kanten
  - positiv is / is-a
  - negativ is-not / is-not-a

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Brachman&Levesque - ch. 10, f.4

## Defeasible Inheritance

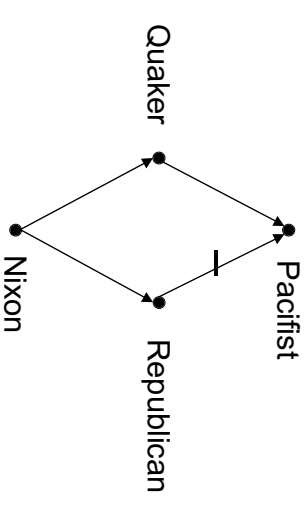
### Gruppen- diskussion

#### Konzeptsystem der Wochentage etc.

- Werktage, Samstag, Sonntag
- Sonn- und Feiertage
- Spezifische Feiertage:
  - Neujahr
  - Karfreitag, Ostern
  - Himmelfahrt
  - Pfingsten
  - 3. Oktober
  - Weihnachten, Sylvester
- Welche Vererbungs-Netze?
- Welche Strategien für Konflikte?

## Defeasible inheritance (2): Nixon Diamond

- Mehrere Vererbungen möglich, die sich aber widersprechen können. (ambiguous net)
- Welche Eigenschaft soll hier zugesprochen werden?
- Welche Strategien für die Entscheidung zwischen den Vererbungs-möglichkeiten?



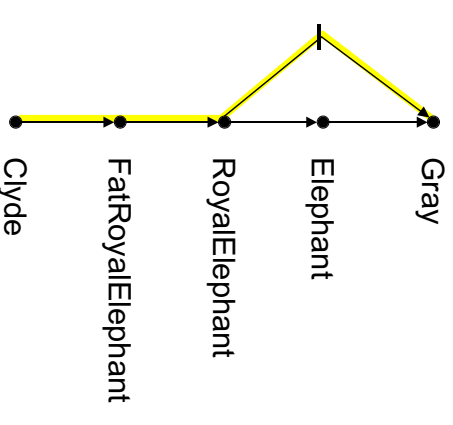
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## Shortest Path Heuristic

#### Defeasible inheritance in DAGs

- links have polarity (positive or negative)
- use **shortest path heuristic** to determine which polarity counts



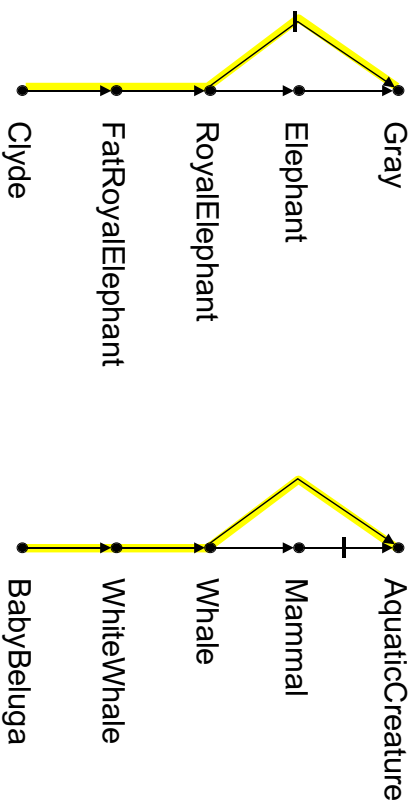
- Zugrundeliegende Idee: Berücksichtige das spezifischste subsumierende Konzept.

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Brachman&Levesque - ch. 10, f.5

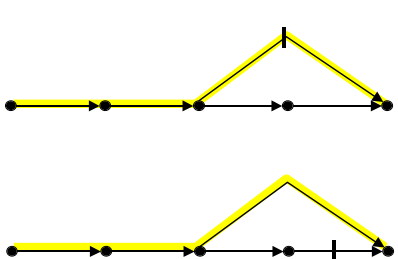
## Shortest Path Heuristic



→ Clyde ist nicht grau.

→ BabyBeluga ist ein im Wasser lebendes Tier.

## Shortest Path Heuristic

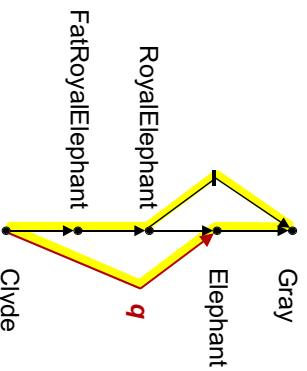


- Pfade im Vererbungsnetzwerk
- fungieren als Argumente
  - stützen Schlussfolgerungen
  - Einige Pfade werden von anderen verhindert (are "preempted")
  - Andere sind zulässig (are "admissible")

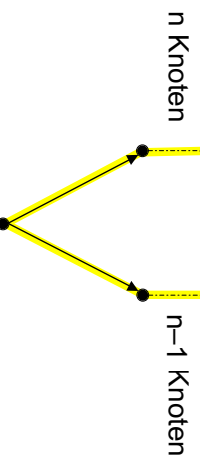
➤ Das Vererbungsproblem ist das Problem, die zulässigen Pfade zu bestimmen.

## Probleme mit der Kürzeste-Wege-Heuristik

Redundante Kanten beeinflussen das Ergebnis, z.B. eine zusätzliche direkte Verbindung, etwa durch einen „gespeicherten Schluss“



Irrelevante Kanten können das Ergebnis beeinflussen, z.B. bei grosser Pfadlänge

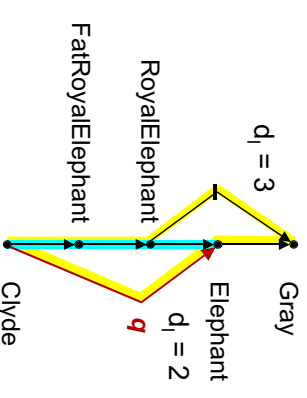


## Inferential distance

Consider "inferential distance": not linear distance, but

- topologically based – a node **a** is nearer to node **b** than to node **c** if there is a path from **a** to **c** through **b**

➤ idea: conclusions from **b** preempt those from **c**



## Inheritance Hierarchy (Formalization L. A. Stein)

An **inheritance hierarchy**  $\Gamma = \langle V, E \rangle$  is a directed, acyclic graph with

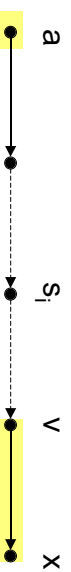
- intended to denote depicted as
- positive edges “(normally) is-a”  $a \rightarrow x$
- and
- negative edges, “(normally) is not-a”  $a \dashrightarrow \neg x$

A sequence of edges is a **path**:

- a **positive path** is a sequence of positive edges  $a \rightarrow \dots \rightarrow x$
- a **negative path** is a sequence of positive edges followed by a single negative edge  $a \rightarrow \dots \rightarrow v \dashrightarrow \neg x$

## Zulässigkeit von Pfaden (Admissibility)

- An edge  $v \dashrightarrow (\neg)x$  is **admissible** in  $\Gamma = \langle V, E \rangle$  w.r.t. **a** if there is a positive path  $a \rightarrow s_1 \dots s_n \rightarrow v$  ( $n \geq 0$ ) in  $E$  and
- each edge in  $a \rightarrow s_1 \dots s_n \rightarrow v$  is admissible in  $\Gamma$  w.r.t. **a** (recursively);
  - no edge in  $a \rightarrow s_1 \dots s_n \rightarrow v$  is **redundant** in  $\Gamma$  w.r.t. **a** (see below);
  - no intermediate node  $a, s_1, \dots, s_n$  is a **preemptor** of  $v \dashrightarrow (\neg)x$  w.r.t. **a** (see below).

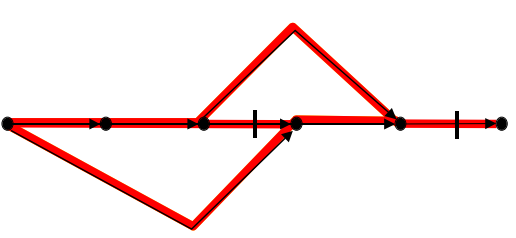


## Vererbungspfade als Argumente

A path (or argument) supports a conclusion:

- $a \rightarrow \dots \rightarrow x$  supports the conclusion  $a \rightarrow x$  ( $a$  is an  $x$ )
- $a \rightarrow \dots \rightarrow v \dashrightarrow \neg x$  supports  $a \dashrightarrow \neg x$  ( $a$  is not an  $x$ )

Note: a conclusion may be supported by many arguments  
However: not all arguments are equally believable...



## Support & Admissibility

- $\Gamma = \langle V, E \rangle$  **supports** a path  $a \rightarrow s_1 \dots s_n \dashrightarrow (\neg)x$  if the corresponding set of edges  $\{a \rightarrow s_1, \dots, s_n \dashrightarrow (\neg)x\}$  is in  $E$ , and it is **admissible**.
- The hierarchy supports a conclusion  $a \rightarrow x$  (or  $a \dashrightarrow \neg x$ ) if it supports some corresponding path

A path is admissible if every edge in it is admissible.

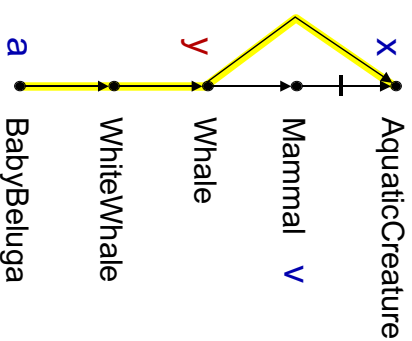
## Preemption (Verhinderung)

A node  $y$  along path

$a \rightarrow \dots \rightarrow y \rightarrow \dots \rightarrow v$  is a **preemptor** of  $v \rightarrow x$  ( $v \rightarrow \neg x$ ) w.r.t  $a$

if  $y \rightarrow \neg x \in E$  ( $y \rightarrow x \in E$ )

Existenz eines Preemptors schliesst Pfade als zulässig aus (Bed. 3).



## Redundanz

A positive edge  $b \rightarrow w$  is redundant in  $\Gamma$  w.r.t. node  $a$  if

there is some positive path

$b \rightarrow t_1 \dots t_m \rightarrow w \in E$  ( $m \geq 1$ ), for which

1. each edge in  $b \rightarrow t_1 \dots t_m$  is admissible in  $\Gamma$  w.r.t.  $a$ ;
2. there are no  $c$  and  $i$  such that  $c \rightarrow \neg t_i$  is admissible in  $\Gamma$  w.r.t.  $a$ ;
3. there is no  $c$  such that  $c \rightarrow \neg w$  is admissible in  $\Gamma$  w.r.t.  $a$ .

The definition for a negative edge

$b \rightarrow \neg w$  is analogous

