

# Wissensrepräsentation

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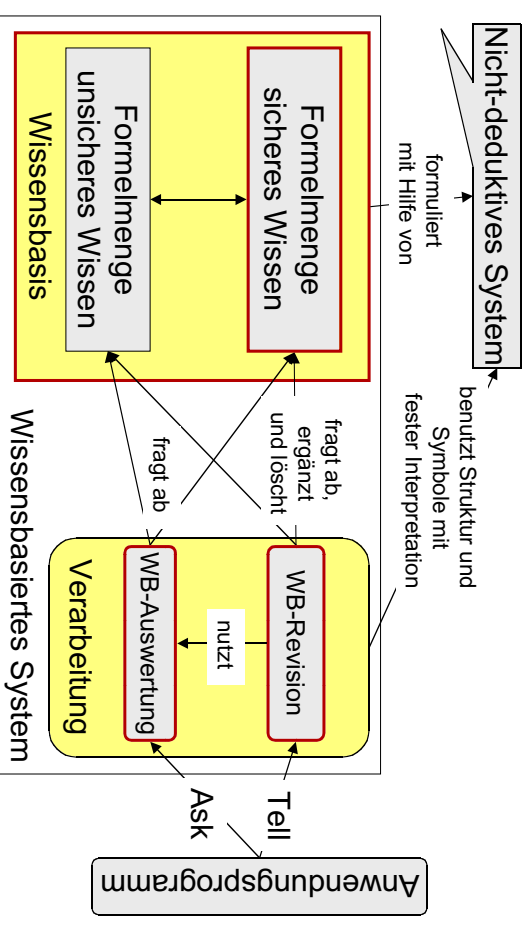
## Sitzung 14: Vererbung / Inheritance (1)

- Vererbung
- Pfadbasiertes Schliessen
- Vererbungsnetze

### Literatur

- D. Gabbay, C. Hogger, and J. Robinson (eds.) (1994). *Handbook of Logic in Artificial Intelligence and Logic Programming, Volume 3: Nonmonotonic Reasoning and Uncertain Reasoning*, Oxford University Press.
- Insbesondere**
- J. Horty. Some direct theories of nonmonotonic inheritance. In D. Gabbay, C. Hogger, and J. Robinson (eds.), Oxford University Press (1994), pp. 111 - 187.
- D. Touretzky. (1986). The mathematics of inheritance systems. Morgan Kaufmann, Los Altos, Ca., 1986.
- L. A. Stein (1992). Resolving Ambiguity in Nonmonotonic Inheritance Hierarchies, *Artificial Intelligence* 55 (2-3): 259-310.

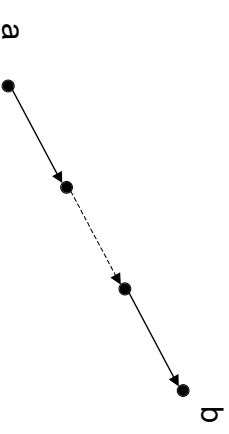
## Wissensbasiertes System mit nicht-deduktiver Wissensbasis



### Inheritance

“Inheritance” is the result of transitivity reasoning over paths

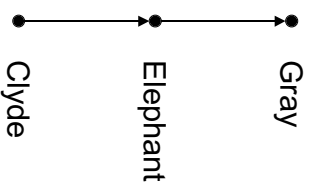
- for strict networks, *modus ponens* in graphical form
- does *a* inherit from *b*?  
 $\equiv$  is *b* in the transitive closure of IS-A (or subsumption) from *a*?



## Pfad-basiertes Schliessen

Focus just on

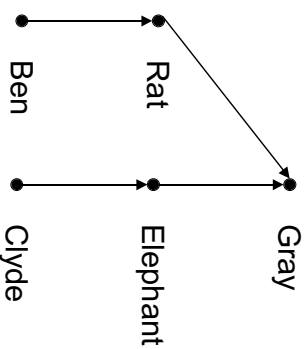
- inheritance and transitivity**
- many interesting considerations in looking just at **where information comes from** in a network representation
- abstract descriptions, and properties into **nodes** in graphs, and just look at **reasoning with paths** and the conclusions they lead us to



## Vererbungsnetze (1)

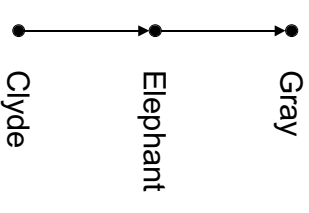
**Strict inheritance in trees**

- as in description logics
- conclusions produced by complete transitive closure on all paths (any traversal procedure will do);
- all reachable nodes are implied



## Pfad-basiertes Schliessen (cont'd)

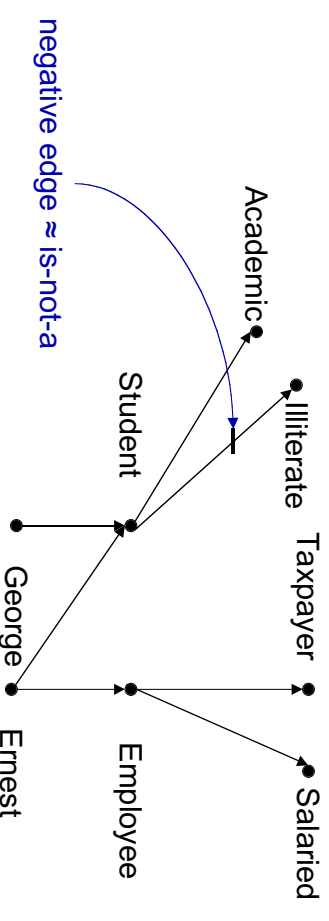
- edges in the network:  
Clyde—Elephant, Elephant—Gray
- paths included in this network:  
{Clyde—Elephant—Gray}, plus edges
- in general, a path is a sequence of 1 or more edges
- conclusions supported by the paths:  
Clyde → Elephant; Elephant → Gray;  
Clyde → Gray



## Vererbungsnetze (2)

**Strict inheritance in DAGs**

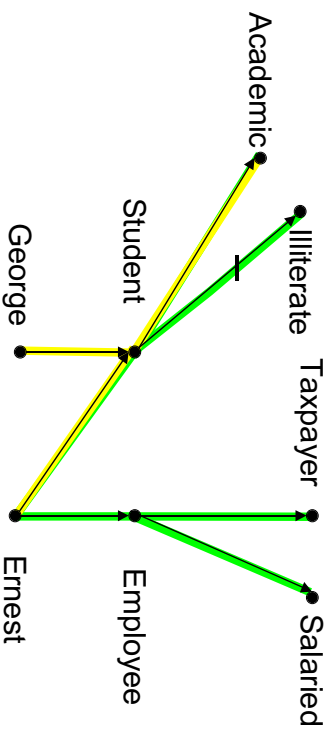
- as in DL's with multiple AND parents
- same as above: all conclusions you can reach by any paths are supported



## Schliessen in Vererbungsnetzen: Upward vs. Downward Reasoning

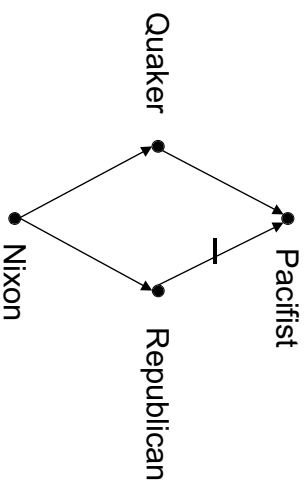
Was ist erschliessbar über:

- Academics
  - Ernest
- ?



## Defeasible inheritance (2): Nixon Diamond

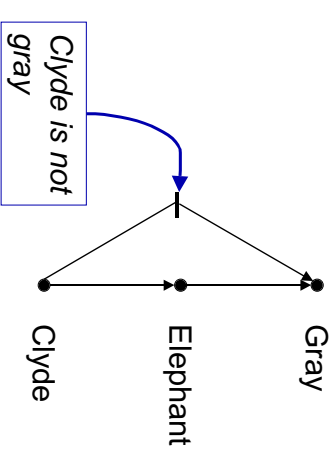
- Mehrere Vererbungen möglich, die sich aber widersprechen können. (ambiguous net)
- Welche Eigenschaft soll hier zugesprochen werden?



- Welche Strategien für die Entscheidung zwischen den Vererbungsmöglichkeiten?

## Defeasible inheritance (1)

- inherited properties do not always hold, and can be **overridden** (defeated)
- conclusions determined by searching upward from “**focus node**” and selecting first version of property you want



- **defeasible**
- **anfechtbar**
- **annullierbar**

- „Polarität“ von Kanten
- positiv is / is-a
- negativ is-not / is-not-a

## Defeasible Inheritance

Aufgabe zur Nachbereitung

### Konzeptsystem der Wochentage etc.

- Werktage, Samstag, Sonntag
- Sonn- und Feiertage
- Spezifische Feiertage:
  - Neujahr
  - Karfreitag, Ostern
  - Himmelfahrt
  - Pfingsten
  - 3. Oktober
  - Weihnachten, Sylvester

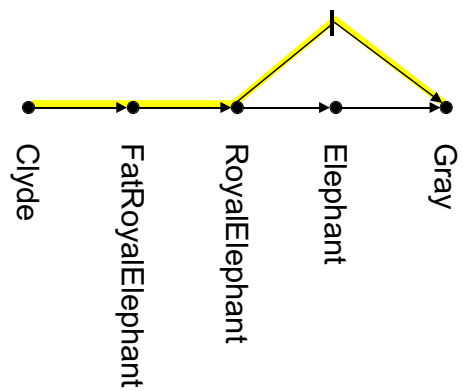
- Welche Vererbungs-Netze?
- Welche Strategien für Konflikte?

## Shortest Path Heuristic

Defeasible inheritance in

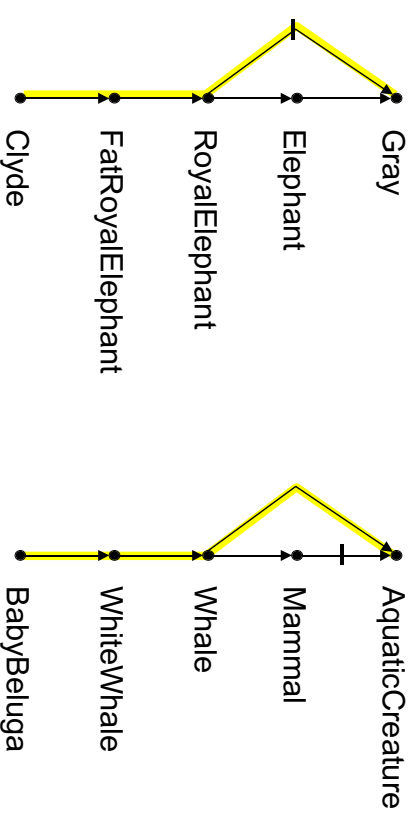
DAGs

- links have polarity (positive or negative)
- use **shortest path heuristic** to determine which polarity counts



- Zugrundeliegende Idee: Berücksichtige das spezifischste subsumierende Konzept.

## Shortest Path Heuristic



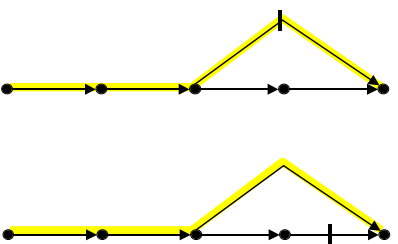
→ Clyde ist nicht grau.

→ BabyBeluga ist ein im Wasser lebendes Tier.

## Shortest Path Heuristic

Pfade im Vererbungsnetzwerk

- fungieren als Argumente
- stützen Schlussfolgerungen
- Einige Pfade werden von anderen verhindert (are "preempted")
- Andere sind zulässig (are "admissible")

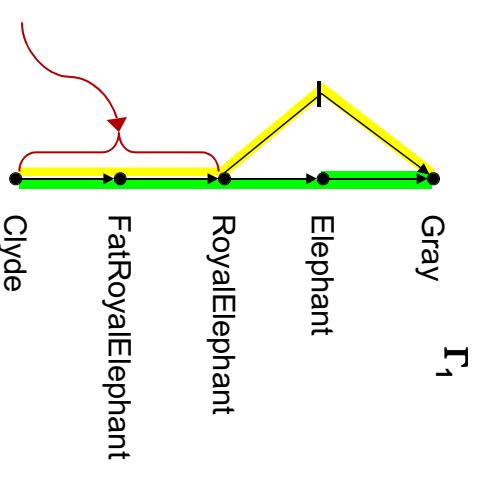


- Das Vererbungsproblem ist das Problem, die zulässigen Pfade zu bestimmen.

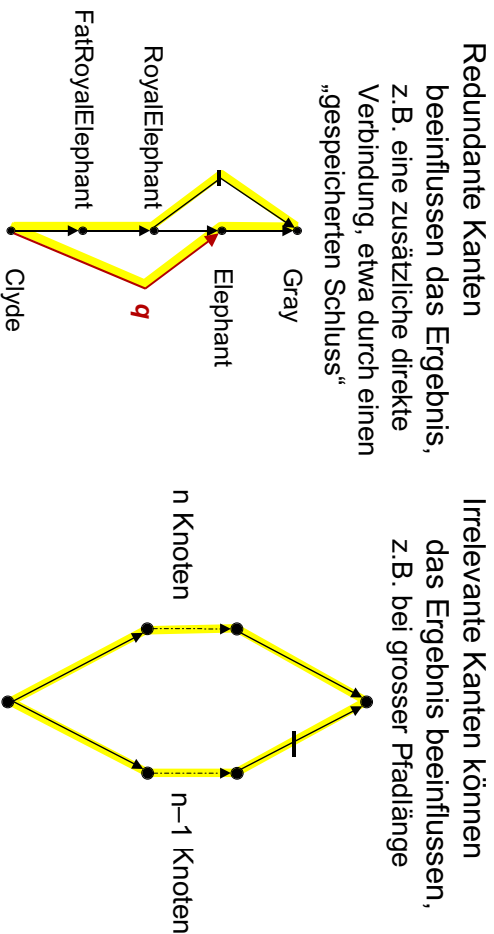
## Pfade und Argumentationen

- Pfade
- werden von einem Knoten ausgehend konstruiert
- sie repräsentieren Argumentationen

Die letzte Kante eines Pfades korrespondiert zu einem Grund, eine bestehende Argumentation fortzusetzen. Verschiedene Argumentationen können gemeinsame Anteile haben.



## Probleme mit der Kürzeste-Wege-Heuristik



## Inheritance Hierarchy (Formalization L. A. Stein, 1992)

An **inheritance hierarchy**  $\Gamma = \langle V, E \rangle$  is a directed, acyclic graph with

- intended to denote “(normally) is-a”
- depicted as  $a \rightarrow x$
- positive edges
- and
- negative edges, “(normally) is not-a”  $a \rightarrow \neg x$

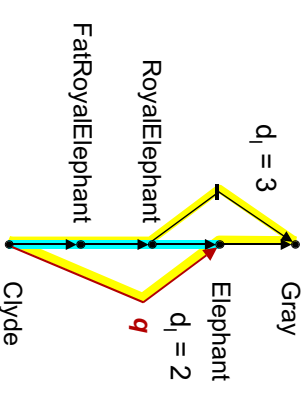
A sequence of edges is a **path**:

- a **positive path** is a sequence of positive edges
- $a \rightarrow \dots \rightarrow x$
- a **negative path** is a sequence of positive edges followed by a single negative edge
- $a \rightarrow \dots \rightarrow v \rightarrow \neg x$

## Inferential distance

Consider “inferential distance”: not linear distance, but

- topologically based – a node **a** is nearer to node **b** than to node **c** if there is a path from **a** to **c** through **b**



➤ idea: conclusions from **b** preempt those from **c**

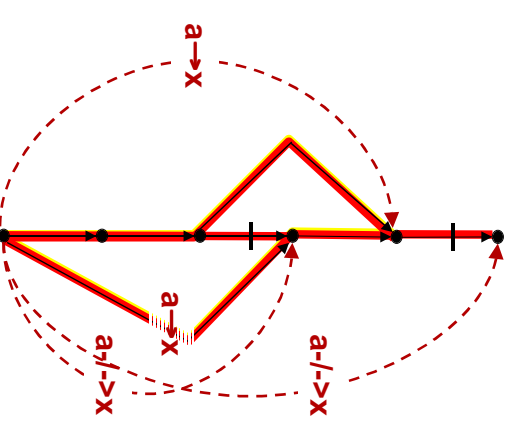
## Vererbungspfade, Argumentationen & Inferenzen

A path (or argument) supports a conclusion:

- $a \rightarrow \dots \rightarrow x$  supports the conclusion  $a \rightarrow x$  (a is an x)
- $a \rightarrow \dots \rightarrow v \rightarrow \neg x$  supports the conclusion  $a \not\rightarrow x$  (a is not an x)

Note:

A conclusion may be supported by many paths  
Different paths can lead to conflicting conclusions.



## Support

$\Gamma = \langle V, E \rangle$  **supports** a path  $a -s_1 \dots -s_n -(\neg)X$  if the corresponding set of edges  $\{a -s_1, \dots, s_n -(\neg)X\}$  is in  $E$ , and it is **admissible**.

$\Gamma \triangleright a -s_1 \dots -s_n -(\neg)X$

- The hierarchy  $\Gamma$  **supports** a conclusion  $a \rightarrow X$  (or  $a -/\rightarrow X$ ) if it supports some corresponding path.

$\Gamma \triangleright a \rightarrow X$  (or  $\Gamma \triangleright a -/\rightarrow X$ )

A path is admissible if every edge in it is admissible.

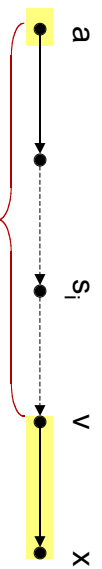
➤ Zu klären ist:

Welche Pfade sind zulässig?

## Zulässigkeit von Pfaden (Admissibility)

An edge  $v -(\neg)X$  is **admissible** in  $\Gamma = \langle V, E \rangle$  w.r.t.  $a$  if there is a positive path  $a -s_1 \dots -s_n -v$  ( $n \geq 0$ ) in  $E$  and

1. each edge in  $a -s_1 \dots -s_n -v$  is admissible in  $\Gamma$  w.r.t.  $a$  (recursively);
2. no edge in  $a -s_1 \dots -s_n -v$  is **redundant** in  $\Gamma$  w.r.t.  $a$  (see below);
3. no intermediate node  $a, s_1, \dots, s_n$  is a **preemptor** of  $v -(\neg)X$  w.r.t.  $a$  (see below).



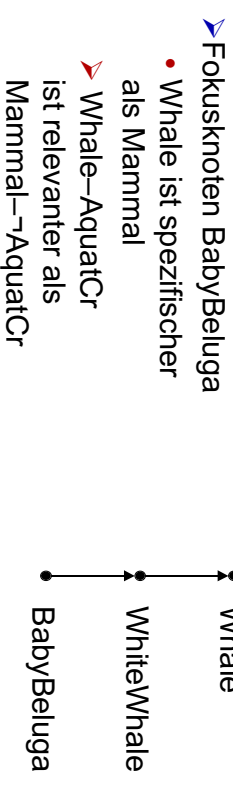
Hier darf nichts vorliegen, was die Kante  $v -(\neg)X$  in ihrer Wirkung innerhalb der Argumentation behindert.

## Spezifizität (specificity):

### Entscheidungen zwischen konfligierenden Pfaden

- Entscheidung fällt auf der Basis von Spezifizität:

- Spezifischere Information sollte mehr Einfluss haben, als weniger spezifische.



## Preemption (Verhinderung)

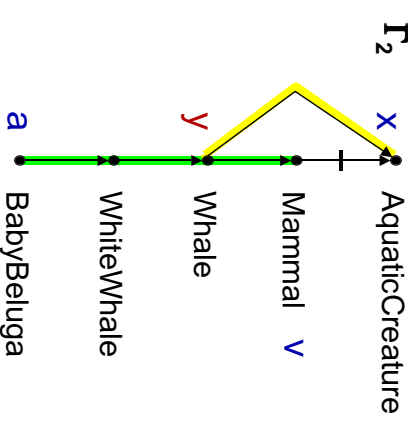
A node  $y$  along path

$a -\dots -y -\dots -v$  is a **preemptor** of  $v -X$  ( $v -\neg X$ ) w.r.t.  $a$

if  $y -\neg X \in E$  ( $y -X \in E$ )

Um die Argumentation  $a -\dots -y -\dots -v$  durch  $v -\neg X$  fortzusetzen bzw. abzuschliessen, muss geprüft werden, ob ein Knoten von  $a -\dots -y -\dots -v$  die weitere Argumentation beeinflusst.

➤ Spezifizität



Preemption wird durch Bedingung 3 Zulässigkeitsdefinition **ausgeschlossen**

## Redundanz

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A positive edge  $b \rightarrow w$  is redundant in  $\Gamma$  w.r.t. node  $a$  if

there is some positive path

$b \rightarrow t_1 \dots t_m \rightarrow w \in E$  ( $m \geq 1$ ), for which

1. each edge in  $b \rightarrow t_1 \dots t_m$  is admissible in  $\Gamma$  w.r.t.  $a$ ;
2. there are no  $c$  and  $i$  such that  $c \rightarrow t_i$  is admissible in  $\Gamma$  w.r.t.  $a$ ;
3. there is no  $c$  such that  $c \rightarrow w$  is admissible in  $\Gamma$  w.r.t.  $a$ .

The definition for a negative edge

$b \rightarrow w$  is analogous

